## Constructing Graphs by Voltage Assignment

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GAP days 2014

## Outline

(1) Voltage Assignment
(2) The Degree-Diameter Problem
(3) GAP Implementation

4 Bibliography

## Voltage graphs

- Given a digraph $G=(V, A)$, and a finite group $\Gamma$, a voltage assignment of $G$ in $\Gamma$ is a function $\alpha: A \rightarrow \Gamma$, that labels the arcs of $G$ with elements of $\Gamma$.
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## Example 1: Voltages in $\mathbb{Z}_{3}$



- Return to GAP session


## Example 1: Collapsing mutually reverse arcs



## Example 2: Cayley graphs (voltages in $\mathbb{Z}_{6}$ )



## Some elementary properties

- The net voltage of a walk $W$ in $G$ is the product of the voltages of every edge in the $W$
- Every cycle $C^{\prime}$ of $G^{\prime}$ corresponds to a closed non-reversing walk $W$ in $G$, with net voltage equal to the identity of $\Gamma$
- The girth of $G^{\prime}$ is equal to the length of the shortest closed non-reversing walk $W$ of $G$, with net voltage equal to the identity
- The local group at vertex $v$ is the group generated by the net voltages of all closed walks based at $v . G^{\prime}$ is connected if and only if the local group at every vertex $v$ is equal to $\Gamma$.


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- Construct the largest possible network (or graph) with a given maximum degree $\Delta$, and a given diameter $D$ (Elspas, 1964)

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## Largest Known Graphs (combinatoricswiki.org)

| $d k$ | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3 | 10 |  |  | 70 | 132 | 196 | 336 | 600 | 1250 |
| 4 |  | 41 | 96 | 364 | 740 | 1320 | 3243 | 7575 | 17703 |
| 5 |  | 72 | 210 | 624 | 2772 | 5516 | 17030 | 57840 | 187056 |
| 6 | 32 | 110 | 390 | 1404 | 7917 | 19383 | 76461 | 307845 | 1253615 |
| 7 | 50 | 168 | 672 | 2756 | 11988 | 52768 | 249660 | 1223050 | 6007230 |
| 8 | 57 | 253 | 1100 | 5060 | 39672 | 131137 | 734820 | 4243100 | 24897161 |
| 9 | 74 | 585 | 1550 | 8200 | 75893 | 279616 | 1686600 | 12123288 | 65866350 |
| 10 | 91 | 650 | 2286 | 13140 | 134690 | 583083 | 4293452 | 27997191 | 201038922 |
| 11 | 104 | 715 | 3200 | 19500 | 156864 | 1001268 | 7442328 | 72933102 | 600380000 |
| 12 | 133 | 786 | 4680 | 29470 | 359772 | 1999500 | 15924326 | 158158875 | 1506252500 |
| 13 | 162 | 851 | 6560 | 40260 | 531440 | 3322080 | 29927790 | 249155760 | 3077200700 |
| 14 | 183 | 916 | 8200 | 57837 | 816294 | 6200460 | 55913932 | 600123780 | 7041746081 |
| 15 | 186 | $\begin{gathered} 1 \\ 215 \end{gathered}$ | $\begin{gathered} 11 \\ 712 \end{gathered}$ | 76518 | $\begin{gathered} 1417 \\ 248 \end{gathered}$ | 8599986 | 90001236 | 1171998164 | $\begin{aligned} & 10012349 \\ & 898 \end{aligned}$ |
| 16 | 198 | $\begin{gathered} 1 \\ 600 \end{gathered}$ | $\begin{aligned} & 14 \\ & 640 \end{aligned}$ | $\begin{aligned} & 132 \\ & 496 \end{aligned}$ | $\begin{gathered} 1771 \\ 560 \end{gathered}$ | $\begin{gathered} 14882 \\ 658 \end{gathered}$ | $\begin{aligned} & 140559 \\ & 416 \end{aligned}$ | 2025125476 | $\begin{aligned} & 12951451 \\ & 931 \end{aligned}$ |
| 17 | 274 | $\begin{aligned} & 1 \\ & 610 \\ & \hline \end{aligned}$ | $\begin{aligned} & 19 \\ & 040 \end{aligned}$ | $\begin{aligned} & 133 \\ & 144 \\ & \hline \end{aligned}$ | $\begin{aligned} & 3217 \\ & 872 \end{aligned}$ | $\begin{aligned} & 18495 \\ & 162 \end{aligned}$ | $\begin{aligned} & 220990 \\ & 700 \end{aligned}$ | 3372648954 | $\begin{aligned} & 15317070 \\ & 720 \end{aligned}$ |
| 18 | 307 | $\left\lvert\, \begin{aligned} & 1 \\ & 620 \end{aligned}\right.$ | $\begin{aligned} & 23 \\ & 800 \end{aligned}$ | $\begin{aligned} & 171 \\ & 828 \\ & \hline \end{aligned}$ | $\begin{aligned} & 4022 \\ & 340 \end{aligned}$ | $\begin{aligned} & 26515 \\ & 120 \\ & \hline \end{aligned}$ | $\begin{aligned} & 323037 \\ & 476 \end{aligned}$ | 5768971167 | $\begin{aligned} & 16659077 \\ & 632 \end{aligned}$ |
| 19 | 338 | $\begin{aligned} & 1 \\ & 638 \end{aligned}$ | $\begin{aligned} & 23 \\ & 970 \end{aligned}$ | $\begin{aligned} & 221 \\ & 676 \end{aligned}$ | $\begin{aligned} & 4024 \\ & 707 \end{aligned}$ | $\begin{aligned} & 39123 \\ & 116 \end{aligned}$ | $\begin{aligned} & 501001 \\ & 000 \end{aligned}$ | 8855 580344 | $\begin{aligned} & 18155097 \\ & 232 \end{aligned}$ |
| 20 | 81 | $\begin{aligned} & 1 \\ & 958 \\ & \hline \end{aligned}$ | $\begin{aligned} & 34 \\ & 952 \\ & \hline \end{aligned}$ | $\begin{aligned} & 281 \\ & 820 \\ & \hline \end{aligned}$ | $\begin{aligned} & 8947 \\ & 848 \end{aligned}$ | $\begin{aligned} & 55625 \\ & 185 \\ & \hline \end{aligned}$ | $\begin{aligned} & 762374 \\ & 779 \end{aligned}$ | $\begin{aligned} & 12951451 \\ & 931 \end{aligned}$ | $\begin{aligned} & 78186295 \\ & 824 \end{aligned}$ |

## Relative contribution of different techniques


(1) Voltage assignment (analytic and computer-based) - $53 \%$
(2) Graph compounding - 15\%
(3) Polarity graphs of generalized polygons-12\%
(9) Other computer-based techniques-9\%
(5) Moore graphs and others - 11\%

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## Random voltage search

Choose a base graph $G$ and a family of groups $\Omega$, and initialize MAX;
Label the arcs of a BFS spanning tree of the base graph $G$ with the identity element;
for every unexplored group $\Gamma$ in $\Omega$ do for i:=1 to MAX do begin generate a random voltage assignment $\alpha$; compute the girth and diameter of $G^{\prime}$;
if diameter $\leq k$ then begin save $\Gamma$ and $\alpha$; break;
end;
end;

## Construction of large general graphs

- The 'less abelian' a group is, the better
- The groups that have been used more extensively are $\mathbb{Z}_{m} \rtimes_{r} \mathbb{Z}_{n}$, $\left(\mathbb{Z}_{m} \times \mathbb{Z}_{m}\right) \rtimes_{r} \mathbb{Z}_{n}$, and $\left(\mathbb{Z}_{m} \rtimes_{r} \mathbb{Z}_{n}\right) \rtimes\left(\mathbb{Z}_{m} \rtimes_{r} \mathbb{Z}_{n}\right)$
- The method is expected to give better results with simple groups, and other non-solvable groups (e.g. perfect groups)
- A group $\Gamma$ is perfect if it equals its commutator (or derived) subgroup $[\Gamma, \Gamma]$. E.g. $S L(2,5)$


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- Package GRAPE, by Leonard Soicher
- Supports directed graphs with loops
- Does not support multiple edges (or multiple loops)
- Does not support edge labels
- Makes heavy use of the automorphism group of the graph
- Calls nauty


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- Group specification
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- Function Lift, to perform the voltage assignment construction
- Function DirectedCayDiameter: Computes the Cayley digraph of a given group 「 that gives the smallest diameter, among all generating sets with given cardinality $k$
- Function UndirectedCayDiameter: the same, for undirected Cayley graphs
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- Euler to/from Pajek's .NET
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- Other possible formais



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## Network layout with Gephi



## Example 1

## Recall Example 1

Gamma:= CyclicGroup (3);
LGamma:= List (Gamma);
vgrafo:= [ [ [ 2 ], [ LGamma[1] ] ],
[ [ 1, 2 ], [ LGamma[1], LGamma[2] ] ] ];
$[$ [4], [5], [6], [1,5], [2,6], [3,4] ]

## Future extensions

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- Improve data structures
- Functions for:
- Adapt functions to different group specifications (permutation groups, fp groups, pc groups, etc.)


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## An alternative data structure

```
Gamma:= CyclicGroup(3);
LGamma:= List(Gamma);
nverts:= 2;
narcs:= 4;
LArcs:= [ [1,2], [2,1], [2,2] ];
alpha:= [ LGamma[1], LGamma[1], LGamma[2] ];
```


## Auxiliary functions (in GAP 3)

- Knuth-Bendix completion procedure for string-rewriting systems: Given a finite monoid presentation, convert it to a complete (confluent) presentation.
- Given a complete presentation, multiply two elements (reduced words)
- Compute the order of a finitely-presented group (monoid) given by a confluent presentation
- Generate confluent presentations for some classes of groups (symmetric groups, alternating groups, finite Coxeter groups, etc.)


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- Compute the order of a finitely-presented group (monoid) given by a confluent presentation
- Generate confluent presentations for some classes of groups (symmetric groups, alternating groups, finite Coxeter groups, etc.)


## Auxiliary functions (in GAP 3)

- Knuth-Bendix completion procedure for string-rewriting systems: Given a finite monoid presentation, convert it to a complete (confluent) presentation.
- Given a complete presentation, multiply two elements (reduced words)
- Compute the order of a finitely-presented group (monoid) given by a confluent presentation
- Generate confluent presentations for some classes of groups (symmetric groups, alternating groups, finite Coxeter groups, etc.)


## Complete presentation of $A_{5}$

```
a:=AbstractGenerator("a");
b:=AbstractGenerator("b") ;
c:=AbstractGenerator("c") ;
```

Gens:=[a,b, c];
Rels: $=$ [[a^3, IdWord], [b^2, IdWord], [ $\left.c^{\wedge} 2, ~ I d W o r d\right]$,
$\left[b * a * b, \quad a^{\wedge} 2 * b * a^{\wedge} 2\right], \quad\left[b * a^{\wedge} 2 * b, \quad a * b * a\right]$,
$\left[c * a, \quad a^{\wedge} 2 * c\right], \quad[c * b * c, \quad b * c * b]$,
$\left.\left[\mathrm{c} * \mathrm{~b} * \mathrm{a} * \mathrm{c}, \quad \mathrm{b} * \mathrm{c} * \mathrm{~b} * \mathrm{a}^{\wedge} 2\right], \quad\left[\mathrm{c} * \mathrm{~b} * \mathrm{a}^{\wedge} 2 * \mathrm{c}, \quad \mathrm{b} * \mathrm{c} * \mathrm{~b} * \mathrm{a}\right]\right] ;$

## GAP-4 translation

```
F:= FreeGroup("a", "b", "C");
a:= F.1;
b:= F.2;
c:= F.3;
IdWord:= One(F);
```

Rels: $=$ [[a^3, IdWord], [b^2, IdWord], [ $c^{\wedge} 2$, IdWord],
$\left[b * a * b, \quad a^{\wedge} 2 * b * a^{\wedge} 2\right], \quad\left[b * a^{\wedge} 2 * b, \quad a * b * a\right]$,
$\left[c * a, \quad a^{\wedge} 2 * c\right], \quad[c * b * c, b * c * b]$,
$\left.\left[c * b * a * c, \quad b * c * b * a^{\wedge} 2\right], \quad[c * b * a \wedge 2 * c, \quad b * c * b * a]\right] ;$

## For more details

B J.L.Gross and T.W.Tucker: Topological Graph Theory. John Wiley \& Sons, 1987.
\& L.H. Soicher: "Computing with graphs and groups" In Topics in Algebraic Graph Theory (L.W. Beineke and R.J. Wilson, eds).
Cambridge Univ. Press, 2004, pp. 250-266.
Q L.H. Soicher: GRAPE Manual.
http://www.gap-system.org/Packages/grape.html.

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Cambridge Univ. Press, 2004, pp. 250-266.
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## END



